Some thoughts on MTA architecture http://dotat.at/writing/mta-arch

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About me

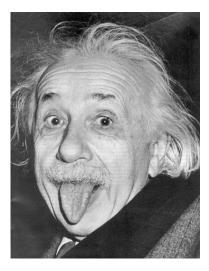
1994 – 1997	$\langle \texttt{fanf2@cam.ac.uk} angle$
1997 – 2000	$\langle \texttt{fanf@demon.net} angle$
2000 - 2001	$\langle \texttt{fanf@covalent.net} \rangle$
2002 – now	$\langle \texttt{fanf2@cam.ac.uk} \rangle$

computer science web server admin Apache httpd coder postmaster

1997	$\langle \texttt{dot@dotat.at} angle$	
1999	$\langle \texttt{fanf@apache.org} angle$	httpd
2002	$\langle \texttt{fanf@FreeBSD.org} \rangle$	unifdef
2004	$\langle \texttt{fanf@exim.org} \rangle$	
2006	$\langle \texttt{fanf@apache.org} \rangle$	${\sf SpamAssassin}$

"Wouldn't it be nice if...?"

- theoretical musings on MTA architecture
- originally a series of postings on my blog, Feb 2006 – March 2007
- there is no code and no likelihood of code



A snapshot of the problem



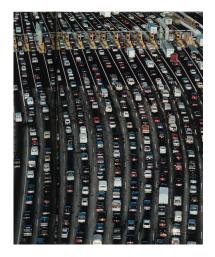
Average email traffic (legitimate and spam): Mar 2005 15

- Mar 2006 20
- Mar 2007 35
- Mar 2008 80
 - all numbers in messages (or rejections) per second

Current traffic classification: 0.5 - 1.5relay attempts known malware 2 - 4blacklisted 60 - 751.5 invalid recipient invalid sender 1.2 SpamAssassin 2 3 legitimate email internal email 2.5

Concurrency

- concurrency requirements grow with spam volumes
- most MTAs use an OS process per connection
- really inefficient!



Waste vs efficiency

- event-driven connection multiplexing
- high-level languages with lightweight threads

better software performance \implies better hardware efficiency



Waste vs efficiency

best use of the available resources ...



Some partial solutions

- SAUCE software against UCE http://www.chiark.greenend.org.uk/~ian/sauce/ (written in Tcl)
- qpsmtpd-async anti-spam smtpd for qmail http://smtpd.develooper.com/ (written in Perl)
- MailChannels Traffic ControlTM http://www.mailchannels.com/products/traffic-control.html

Address verification

- most verifications are for messages that will be rejected
- email address routeing can be arbitrarily complicated so verification can be too!
- concurrency useful for multi-recipient messages as well as multiple messages

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Avoid bouncing

- reject unwanted email as early as possible
- try hard not to accept and bounce
- reduce spam backscatter & forwarded spam
- avoid wasting your MTA's resources



How email addresses are routed

- DNS MX/A/AAAA
- flat files text or cdb
 - aliases
 - mailertable
 - virtusertable
- LDAP "laser" schema
- SQL databases



User-defined filtering

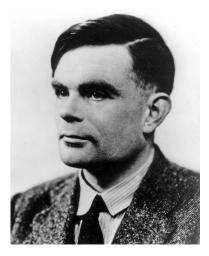
Sieve — RFC 5228

- address validity can be conditional on the sender's address
- selective sub-address validity, e.g. fanf9+subaddress@hermes.cam.ac.uk



Routeing with regular expressions

- try to match address against a series of regular expressions
- when one matches, replace address with corresponding result
- interpolate captured subexpressions
- route resulting address, repeating regsub if necessary



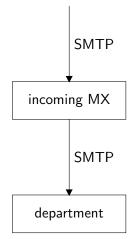
Verification: you're doing it wrong!

Postfix local_recipients_map



Verifying relayed addresses





Verification: you're doing it wrong!

- copy table of valid recipients from department to MX
- configure MX to query department's LDAP directory



Call-forward recipient verification

. . .

```
220 mx.cam.ac.uk
HELO dotat.at
250 Hello
MATL FROM: <dot@dotat.at>
250 OK
RCPT TO:<?@cl.cam.ac.uk> 220 mta.cl.cam.ac.uk
                           HELO mx.cam.ac.uk
                           250 Hello
                           MAIL FROM: <dot@dotat.at>
                           250 OK
                           RCPT TO:<?@cl.cam.ac.uk>
                           550 Unknown user
550 Unknown user
                           QUIT
RSET
                           221 Goodbye
```

Content scanning



- ▶ anti-spam
- anti-phishing
- ► anti-virus
- ► lots of CPU
- Iots of memory

Content scanning goals

- decouple scanner from client concurrency & speed
- do not require entire message to be buffered in RAM
- avoid temporary on-disk buffers
- security boundary between content scanner(s) and MTA



Data callout

- use the normal local delivery mechanism
- efficiently transfer a file from the queue to a program
- cross security boundaries
- control concurrency and smooth load spikes



Queue layout

- MTAs typicall scatter messages all over the disk
- often separate files for envelopes and contents
- this makes queue runs particularly expensive

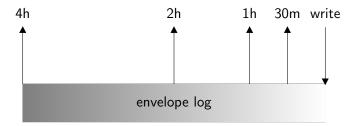


Log-structured queue

- write all metadata sequentially to one file
- queue runners read file sequentially
- updated envelopes also appended to the file
- queue runners act as garbage collectors
- size of log bounded by retry interval



Log-structured queue





Architectural principles

- lightweight concurrency througout the system
- load smoothing / scheduling of scarce resources
 - database connections, content scanners
- address routeing is verification
- content scanning is a data call-forward
- a log-structured queue minimizes disk seeks



That's all, folks!



- slides and notes available online: http://dotat.at/writing/mta-arch
- any questions?